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# Boundedness of oriented walks generated by substitutions 

par F.M. DEKKING et Z.-Y. WEN

Résumé. Soit $x=x_{0} x_{1} \ldots$ un point fixe de la substitution sur l'alphabet $\{a, b\}$, et soit $U_{a}=\left(\begin{array}{rr}-1 & -1 \\ 0 & 1\end{array}\right)$ et $U_{b}=\left(\begin{array}{ll}1 & 1 \\ 0 & 1\end{array}\right)$. On donne une classification complète des substitutions $\sigma:\{a, b\}^{*} \rightarrow\{a, b\}^{*}$ selon que la suite de matrices $\left(U_{x_{0}} U_{x_{1}} \ldots U_{x_{n}}\right)_{n=0}^{\infty}$ est bornée ou non. Cela correspond au fait que les chemins orientés engendrés par les substitutions sont bornés ou non.

Abstract. Let $x=x_{0} x_{1} \ldots$ be a fixed point of a substitution on the alphabet $\{a, b\}$, and let $U_{a}=\left(\begin{array}{rr}-1 & -1 \\ 0 & 1\end{array}\right)$ and $U_{b}=\left(\begin{array}{ll}1 & 1 \\ 0 & 1\end{array}\right)$. We give a complete classification of the substitutions $\sigma:\{a, b\}^{*} \rightarrow\{a, b\}^{*}$ according to whether the sequence of matrices $\left(U_{x_{0}} U_{x_{1}} \ldots U_{x_{n}}\right)_{n=0}^{\infty}$ is bounded or unbounded. This corresponds to the boundedness or unboundedness of the oriented walks generated by the substitutions.

## 1. Introduction

Let $A$ be the alphabet $\{a, b\}$, and let $x=x_{0} x_{1} \ldots$ be an infinite sequence over $A$. Any such sequence generates an oriented walk $\left(S_{N}\right)=\left(S_{N, f}(x)\right)$ on the integers by the following rules:

$$
\begin{align*}
S_{-1} & =-1, \quad S_{0}=0  \tag{1}\\
S_{N+1} & =\left\{\begin{array}{lll}
S_{N-1} & \text { if } & x_{N}=a \\
2 S_{N}-S_{N-1} & \text { if } & x_{N}=b
\end{array}\right.
\end{align*}
$$

In other words: we move one step in the same direction if $x_{N}=b$, and one step in the reversed direction if $x_{N}=a$. Another way to describe $\left(S_{N}\right)_{N=0}^{\infty}$ is by introducing the matrices

$$
U_{a}=\left(\begin{array}{rr}
-1 & -1 \\
0 & 1
\end{array}\right), U_{b}=\left(\begin{array}{ll}
1 & 1 \\
0 & 1
\end{array}\right)
$$

then

$$
S_{N}(x)=\left(\begin{array}{ll}
1 & 0
\end{array}\right) U_{x_{0}} U_{x_{1}} \ldots U_{x_{N-1}}\binom{0}{1}
$$

[^0]Manuscrit reçu le 3 juin 1995.

In the probability literature $\left(S_{N}(x)\right)$ is also known as persistent random walk or correlated random walk if $x$ is obtained according to a product measure on $A^{\mathbb{N}}$. Here we shall consider the case where the sequence $x$ is a fixed point of a primitive substitution $\sigma$ on $\{a, b\}$.
Non-oriented walks $\left(S_{N, f}(x)\right)$ on R are defined by $S_{0}(x)=0$ and

$$
S_{N, f}(x)=\sum_{k=0}^{N-1} f\left(x_{k}\right) \quad \text { for } N \geq 1
$$

where $f: A \rightarrow \mathbf{R}$, and $A$ is now an arbitrary finite set. (It is convenient to extend $f$ homomorphically to $A^{*}=\cup_{k \geq 0} A^{k}$, i.e. $f\left(w_{1} \ldots w_{k}\right)=f\left(w_{1}\right)+\cdots+f\left(w_{k}\right)$ for $\left.w_{1} \ldots w_{k} \in A^{*}\right)$.
Non-oriented walks with $x$ a fixed point of a substitution have been studied in [2], [3], [5], [6], [7], [10]. Here [10] contains a rather complete analysis of the behaviour of $\left(S_{N, f}(x)\right)$ for two letter alphabets $A=\{a, b\}$.
It follows from a general result in [4] that by enlarging the alphabet $A$ oriented walks generated by substitutions may be viewed as non-oriented walks generated by substitutions. Hence it might look as if the main result of [5],[6] - which admits alphabets of arbitrary sizes - would answer all questions on the two symbol oriented walk. Their result is that as $N \rightarrow \infty$
(2) $\quad S_{N, f}(x)=(v, f) N+\left(\log _{\theta}(N)\right)^{\alpha} N^{\log _{\theta}\left(\theta_{2}\right)} F(N)+o\left((\log (N))^{\alpha} N^{\log _{\theta}\left(\theta_{2}\right)}\right)$
where $\theta$ is the Perron-Frobenius eigenvalue of the matrix $M_{\sigma}$ of the substitution $\sigma$ (with entries $m_{s t}=|\sigma(t)|_{s}=$ number of occurrences of the symbol $s$ in the word $\sigma(t)$ ), $\theta_{2}$ the second largest (in absolute value) eigenvalue of $M_{\sigma}$ (which is required to be unique and larger than 1) and $v$ is the vector satisfying $M_{\sigma} v=\theta v$ and $\sum_{s \in A} v_{s}=1$. Furthermore $\alpha+1$ is the order of $\theta_{2}$ in the minimal polynomial of $M$, and $F:[1, \infty) \rightarrow \mathbf{R}$ is a bounded continuous function which satisfies a self-similarity property:

$$
F(\theta x)=F(x) \quad x \geq 1
$$

However, even such a simple property as boundedness of $\left(S_{N, f}(x)\right)_{N \geq 0}$ can often not be resolved with (2). Let us take for example $A=\{a, b, c\}, f(a)=f(b)=$ $+1, f(c)=-1$ and $\sigma$ such that the matrix of $\sigma$ equals

$$
M_{\sigma}=\left(\begin{array}{lll}
3 & 1 & 2 \\
1 & 3 & 2 \\
4 & 4 & 4
\end{array}\right)
$$

We quickly see that $M_{\sigma}$ has eigenvalues $0, \theta_{2}=2$ and $\theta=8$, and that the Perron Frobenius eigenvector $v=(1,1,2)^{T}$ satisfies $(v, f)=0$. We obtain from (2) that as $N \rightarrow \infty$

$$
\begin{equation*}
S_{N, f}(x)=N^{1 / 3} F(N)+o\left(N^{1 / 3}\right) \tag{3}
\end{equation*}
$$

But $f(\sigma a)=f(\sigma b)=f(\sigma c)=0$, so $S_{8 N, f}(x)=0$ for all $N$, and $\left(S_{N, f}(x)\right)$ is bounded, in spite of the behaviour suggested by (3). (Of course (3) and the self-similarity property of $F$ imply that $F \equiv 0$.)

The goal of this study is to determine for any substitution $\sigma$ on $\{a, b\}$ whether the oriented walk in (1) will be bounded or not.

Although not explicitly formulated, the analysis of the oriented two symbol case in [10] heavily relies on the fact that a substitution $\sigma$ on a two symbol alphabet with $(v, f)=0$ automatically admits a representation with the same $f$ in $\mathbf{R}$ (terminology from [1]), i.e., there exists $\lambda \in \mathbb{R}$ such that $f(\sigma(s))=\lambda f(s)$ for $s=a, b$. (In [9] such $\sigma$ are called geometric, see also [8]). However this is no longer true for larger alphabets, and this is the main reason that our solution to the boundedness problem is rather delicate.

## 2. Four types of substitutions

Let $\sigma$ be a substitution on $\{a, b\}$ such that the first letter of $\sigma(a)$ is $a$, and let $u$ be the fixed point of $\sigma$ with $u_{0}=a$. Let

$$
M_{\sigma}=\left(\begin{array}{ll}
|\sigma a|_{a} & |\sigma b|_{a} \\
|\sigma a|_{b} & |\sigma b|_{b}
\end{array}\right)
$$

be the matrix of $\sigma$. Here, as usual, $|v|_{w}$ denotes the number of occurrences of a word $w$ in a word $v$. It appears that the question of boundedness of ( $\left.S_{N}(u)\right)$ depends crucially on the entries of $M_{\sigma}$ reduced modulo two. Let $\bar{M}_{\sigma}$ be this matrix. Then there are $2^{4}=16$ of these matrices possible. However, since $\sigma, \sigma^{2}$ and $\sigma^{3}$ all generate the same fixed point $u$, we only have to consider four types, namely

$$
\text { (I) } \bar{M}_{\sigma}=\left(\begin{array}{ll}
0 & 0 \\
\ldots
\end{array}\right), \text { (II) } \bar{M}_{\sigma}=\left(\begin{array}{ll}
1 & 1 \\
0 & 0
\end{array}\right), \text { (III) } \bar{M}_{\sigma}=\left(\begin{array}{cc}
1 & 0 \\
\ldots
\end{array}\right), \text { (IV) } \bar{M}_{\sigma}=\left(\begin{array}{ll}
0 & 1 \\
0 & 1
\end{array}\right) .
$$

Here the dots indicate that the corresponding entries are either 0 or 1 . The four cases cover respectively $6,1,8$ and 1 of the 16 possibilities. For example the Fibonacci substitution $a \rightarrow a b, b \rightarrow a$ belongs to Type III since $\left(\begin{array}{ll}1 & 1 \\ 1 & 0\end{array}\right)^{3} \equiv\left(\begin{array}{ll}1 & 0 \\ 0 & 1\end{array}\right)$.

## 3. Type I substitutions

Here $\bar{M}_{\sigma}=\binom{0}{0}$. The essential feature of this case is that the number of $a$ 's in both $\sigma(a)$ and $\sigma(b)$ being even, the orientation at the beginning is the same as at the end of these words. Hence if we consider $\sigma(a)$ and $\sigma(b)$ as new symbols we can obtain a non-oriented walk which behaves very much as the original oriented walk. Formally, define the homomorphism (w.r.t. concatenation)

$$
\varphi:\{\sigma(a), \sigma(b)\}^{*} \rightarrow\{\alpha, \beta\}^{*}
$$

by $\varphi(\sigma(a))=\alpha, \varphi\left((\sigma(b))=\beta\right.$. Then define a substitution $\hat{\sigma}$ on $\{\alpha, \beta\}^{*}$ by

$$
\hat{\sigma}(\alpha)=\varphi\left(\sigma^{2} a\right), \quad \hat{\sigma}(\beta)=\varphi\left(\sigma^{2} b\right)
$$

(actually $\sigma=\hat{\sigma}$ but for a change of alphabet!), and define $\hat{f}:\{\alpha, \beta\} \rightarrow \mathbf{R}$ by

$$
\hat{f}(\alpha)=S_{\ell_{a}}(\sigma(a)), \quad \hat{f}(\beta)=S_{\ell_{b}}(\sigma(b))
$$

where $\ell_{a}=|\sigma(a)|$ (the length of $\sigma(a)$ ), and $\ell_{b}=|\sigma(b)|$.
Let $\hat{u}$ be the fixed point of $\hat{\sigma}$ with $\hat{u}_{0}=\alpha$. Then the non-oriented walk $\left(\hat{S}_{N, \hat{f}}(\hat{u})\right.$ ) visits a subsequence of the original oriented walk $\left(S_{N}(u)\right)$, the time instants not being further apart than $\max \left(\ell_{a}, \ell_{b}\right)$. Hence boundedness of $\left(S_{N}(u)\right)$ is equivalent to boundedness of $\left(\hat{S}_{N, \hat{f}}(\hat{u})\right)$. The latter can be easily resolved with Theorem 1.27 of [10].
Example. Let $\sigma$ be the Prouhet-Thue-Morse substitution $\sigma(a)=a b b a, \sigma(b)=$ $b a a b$. Then $f(\alpha)=-2, f(\beta)=2$ and $\hat{\sigma}(\alpha)=\alpha \beta \beta \alpha, \hat{\sigma}(\beta)=\beta \alpha \alpha \beta$. It is easy to see that $\left(\hat{S}_{N, \hat{f}}(\hat{u})\right) \in\{-2,0,2\}$, so the original oriented walk is also bounded (actually it is confined to the set $\{-3,-2,-1,0,1,2\}$ ).

## 4. An equivalence relation

We call words $v, w \in A^{*}=\cup_{k \geq 0} A^{k}$ of length $n$ and length $m$ equivalent, and denote this by $v \sim w$ if

$$
S_{n-1}(v)=S_{m-1}(w) \text { and } S_{n}(v)=S_{m}(w)
$$

i.e., the associated oriented walks end at the same integer with the same orientation. In terms of the matrices $U_{a}$ and $U_{b}$ introduced in Section 1 we have $v \sim w$ iff (1 0) $U_{v}=(10) U_{w}$ (here $U_{w_{1} w_{2} \ldots w_{k}}=U_{w_{1}} U_{w_{2}} \ldots U_{w_{k}}$ if $w \in A^{k}$ ). Note that concatenation preserves equivalence. We denote the empty word by $\epsilon$. Typical examples are

$$
a^{2} \sim \epsilon \quad, \quad a b a b \sim \epsilon
$$

Since the orientation changes iff an $a$ occurs we have
Lemma 1. If $v \sim w$, then $|v|_{a} \equiv|w|_{a}$ modulo 2.
The following lemma is important in the analysis of Type II and IV.
Lemma 2. For all $w \in A^{*}$ there exist $r, \ell \in\{0,1\}$ and $n \in \mathbb{N}$ such that $w \sim$ $a^{\ell} b^{n} a^{r}$.
Proof. Apply $a^{2} \sim \epsilon$ until $w \sim a^{\ell} b^{n_{1}} a b^{n_{2}} a \ldots a b^{n_{k}} a^{r}$ for some $k$. Then apply $b a b \sim a \min \left(n_{k-1}, n_{k}\right)$ times on $b^{n_{k-1}} a b^{n_{k}}$. The result is that $w$ is equivalent to a word of the form above with $k$ one smaller. The lemma then follows by induction.

Lemma 3 (Squaring Lemma). If $|w|_{a}$ is odd then $w^{2} \sim \epsilon$.

Proof. By Lemma 2, $w^{2} \sim a^{\ell} b^{n} a^{r+\ell} b^{n} a^{r}$, and by Lemma 1, $r+\ell$ is odd. So $w^{2} \sim a^{\ell} b^{n} a b^{n} a^{r}$. But since $b^{n} a b^{n} \sim a$, we obtain $w^{2} \sim a^{\ell+r+1} \sim \epsilon$.

Warning: note that in general $u \sim v$ does not imply $\sigma(u) \sim \sigma(v)$.

## 5. Type II substitutions

Here $\bar{M}_{\sigma}=\left(\begin{array}{ll}1 & 1 \\ 0 & 0\end{array}\right)$. Note that also $\bar{M}_{\sigma}^{n}=\left(\begin{array}{ll}1 & 1 \\ 0 & 0\end{array}\right)$. Hence the Squaring Lemma implies

$$
\begin{equation*}
\sigma^{n}\left(a^{2}\right) \sim \epsilon \quad, \quad \sigma^{n}\left(b^{2}\right) \sim \epsilon \quad \text { for all } n \geq 1 \tag{4}
\end{equation*}
$$

Proposition 1. If $\sigma$ is of Type II, then $\left(S_{N}(u)\right)$ is bounded iff $\sigma^{2}(a b) \sim \epsilon$.
Proof. Note first that $\sigma^{2}(a b) \sim \epsilon$ iff $\sigma^{2}(b a) \sim \epsilon$. Since if for example $\sigma^{2}(a b) \sim \epsilon$, then

$$
\sigma^{2}(b a) \sim \sigma^{2}\left(a^{2}\right) \sigma^{2}(b a)=\sigma^{2}(a) \sigma^{2}(a b) \sigma^{2}(a) \sim \sigma^{2}\left(a^{2}\right) \sim \epsilon
$$

$\Leftrightarrow)$ Now $u=\sigma^{2}(u)=\sigma^{2}\left(u_{0}\right) \sigma^{2}\left(u_{1}\right) \ldots$, where each $\sigma^{2}\left(u_{2 k} u_{2 k+1}\right) \sim \epsilon$.
This obviously implies that $\left(S_{N}(u)\right)$ is bounded.
$\Rightarrow)$ We will show that $\sigma^{2}(a b) \nsim \varepsilon$ implies that $\left(S_{N}(u)\right)$ is unbounded. Because of (4) any word $\sigma^{n}(w)$ is equivalent to one of the following words for some $k \geq 0$

$$
\left[\sigma^{n}(a b)\right]^{k},\left[\sigma^{n}(b a)\right]^{k}, \sigma^{n}(b)\left[\sigma^{n}(a b)\right]^{k} \text { or } \sigma^{n}(a)\left[\sigma^{n}(b a)\right]^{k}
$$

Now we take for $w$ the word $\sigma(a b)$. Since the numbers of $a$ 's and $b$ 's in $\sigma(a b)$ are both even, only the first two possibilities above remain, and moreover, $k$ is even. Let us consider the first possibility, i.e., $\sigma^{n+1}(a b) \sim\left[\sigma^{n}(a b)\right]^{k}$. Then also $\sigma^{n}(a b)=\sigma^{n-1}(\sigma(a b))=\left[\sigma^{n-1}(a b)\right]^{k}$, hence

$$
\sigma^{n+1}(a b) \sim\left[\sigma^{n-1}(a b)\right]^{k^{2}} .
$$

Continuing in this fashion we obtain

$$
\begin{equation*}
\sigma^{n+1}(a b) \sim\left[\sigma^{2}(a b)\right]^{k^{n-1}} \tag{5}
\end{equation*}
$$

Since we assume that $\sigma^{2}(a b) \nsim \varepsilon$, and since $\sigma^{2}(a b) \sim[\sigma(a b)]^{k}$, we have $k>0$, so $k \geq 2$. Since $a b$, and hence $\sigma^{n+1}(a b)$ has to occur (5) implies that $\left(S_{N}(u)\right)$ is unbounded, because $\sigma^{2}(a b)$ contains an even number of $a$ 's which implies that the walk corresponding to $\sigma^{2}(a b)$ does not change orientation. In case $\sigma^{n+1}(a b) \sim$ $\left[\sigma^{n}(b a)\right]^{k}$ the same argument applies with $a$ and $b$ interchanged.
Example. We consider 2 substitutions with matrix $\left(\begin{array}{ll}3 & 1 \\ 2 & 2\end{array}\right)$.
A. Let $\sigma(a)=a a b a b, \sigma(b)=b b a$.

Then $\sigma^{2}(a b) \sim[\sigma(b a)]^{2} \sim b^{4}$, so the walk is unbounded.
B. Let $\sigma(a)=a b b a a, \sigma(b)=a b b$.

Then $\sigma^{2}(a b) \sim \epsilon$, so the walk is bounded.
(In case B also $\sigma(a b) \sim \epsilon$. The substitution given by $\sigma(a)=a b b, \sigma(b)=a$ provides an example where $\sigma^{2}(a b) \sim \epsilon$, but $\sigma(a b) \nsim \epsilon$ ).

## 6. Type III substitutions

Here $\bar{M}_{\sigma}=\left(\begin{array}{ll}1 \\ 0 \\ 0\end{array}\right)$. Now the orientation has not changed after occurrence of $\sigma(b)$. The idea is then to keep track of the parity of the number of $a$ 's that have occurred until occurrence of $u_{n}$ in $u$, and obtain $\left(S_{N}(u)\right)$ as non-oriented walk $\left(S_{N, \dot{f}}(\dot{u})\right)$. To this end we consider a four symbol alphabet $\dot{A}=\left\{a^{+}, a^{-}, b^{+}, b^{-}\right\}$ with a substitution $\dot{\sigma}$ with fixed point $\dot{u}$. E.g. $\dot{u}_{n}=a^{+}$will mean that $u_{n}=a$ and that an even number of $a$ 's have occurred in $u_{0} \ldots u_{n-1}$. The substitution $\dot{\sigma}$ is defined by exponentiating the symbols $a$ and $b$ of $\sigma(a)$ and $\sigma(b)$, by +'s and -'s according to the rules: (i) the first symbol obtains a + , (ii) if a symbol follows an $a$ the exponent is reversed if it follows a $b$ it remains equal to that of its predecessor. The $\dot{\sigma}\left(a^{-}\right)$and $\dot{\sigma}\left(b^{-}\right)$are obtained by reversing the signs in $\dot{\sigma}\left(a^{+}\right)$, respectively $\dot{\sigma}\left(b^{+}\right)$. Now we define $\dot{f}: A \rightarrow \mathbf{R}$ by

$$
\dot{f}\left(a^{+}\right)=\dot{f}\left(b^{+}\right)=1, \quad \dot{f}\left(a^{-}\right)=\dot{f}\left(b^{-}\right)=-1
$$

Then it maybe verified (this is a special case of the construction in [4]) that for $N=-1,0,1, \ldots$

$$
S_{N}(u)=S_{N+1, \dot{f}}(\dot{u})-1
$$

where $\dot{u}$ is the fixed point of $\dot{\sigma}$ with $\dot{u}_{0}=a^{+}$.
Example. Let $\sigma$ be given by $\sigma(a)=a a b a b, \sigma(b)=a b a b b$. This induces a substitution $\dot{\sigma}$ on the alphabet $\left\{a^{+}, a^{-}, b^{+}, b^{-}\right\}$by

$$
\begin{array}{ccc}
\dot{\sigma}\left(a^{+}\right)=a^{+} a^{-} b^{+} a^{+} b^{-}, & & \dot{\sigma}\left(a^{-}\right)=a^{-} a^{+} b^{-} a^{-} b^{+} \\
\dot{\sigma}\left(b^{+}\right)=a^{+} b^{-} a^{-} b^{+} b^{+}, & & \dot{\sigma}\left(b^{-}\right)=a^{-} b^{+} a^{+} b^{-} b^{-}
\end{array}
$$

By definition sign changes ( $a^{+}$followed by $a^{-}$or $b^{-}$, etc.) occur and only occur in the word $\dot{\sigma}(\cdot)$ directly following a symbol $a^{+}$or $a^{-}$. Since $\bar{M}_{\sigma}=\left({ }^{1}{ }^{0}.\right)$, this implies that in $\dot{\sigma}\left(a^{+}\right)$there is exactly one more $a^{+}$, say $y+1$, than $a^{-}$(note that $\dot{\sigma}\left(a^{+}\right)$always starts with $\left.a^{+}\right)$, and in $\dot{\sigma}\left(b^{+}\right)$there are equal numbers of $a^{+}$ and $a^{-}$, say $x$. Moreover, one obtains $\dot{\sigma}\left(a^{-}\right)$from $\dot{\sigma}\left(a^{+}\right)$by reversing signs, and similarly for $\dot{\sigma}\left(b^{-}\right)$. Combining these constraints, we see that the matrix $M_{\dot{\sigma}}$ of $\dot{\sigma}$ has the form

$$
M_{\dot{\sigma}}=\left(\begin{array}{cccc}
y+1 & y & x & x \\
y & y+1 & x & x \\
s & t & p & q \\
t & s & q & p
\end{array}\right)
$$

where all entries are non-negative. It turns out that the crucial parameters are $\alpha$ and $\beta$ defined by

$$
\begin{aligned}
& \alpha:=s-t=\left|\dot{\sigma}\left(a^{+}\right)\right|_{b^{+}}-\left|\dot{\sigma}\left(a^{+}\right)\right|_{b^{-}} \\
& \beta:=p-q=\left|\dot{\sigma}\left(b^{+}\right)\right|_{b^{+}}-\left|\dot{\sigma}\left(b^{+}\right)\right|_{b^{-}} .
\end{aligned}
$$

Proposition 2. If $\sigma$ if of Type III, then $\left(S_{N}(u)\right)$ is bounded iff $\beta=0$, or if there exist even numbers $m_{a}$ and $m_{b}$ such that

$$
\sigma(a)=(a b)^{m_{a}} a, \quad \sigma(b)=(b a)^{m_{b}} b .
$$

Proof. For real numbers $K, L$ let

$$
w_{K, L}=(K,-K, L,-L)
$$

Then

$$
\begin{equation*}
w_{K, L} M_{\dot{\sigma}}=(K+\alpha L,-K-\alpha L, \beta L,-\beta L) \tag{6}
\end{equation*}
$$

Therefore one has for all $n \geq 1$

$$
\begin{equation*}
w_{K, L} M_{\dot{\sigma}}^{n}=\left(K+\alpha L\left(1+\beta+\cdots+\beta^{n-1}\right), \ldots, \beta^{n} L,-\beta^{n} L\right) \tag{7}
\end{equation*}
$$

Note that

$$
\begin{equation*}
\left(\dot{f}\left(a^{+}\right), \dot{f}\left(a^{-}\right), \dot{f}\left(b^{+}\right), \dot{f}\left(b^{-}\right)\right)=(1,-1,1,-1)=w_{1,1} \tag{8}
\end{equation*}
$$

So, if $\ell_{n}=\left|\dot{\sigma}^{n}\left(a^{+}\right)\right|$, then

$$
S_{\ell_{n}, \dot{f}}(\dot{u})=\sum_{k=0}^{\ell_{n}-1} \dot{f}\left(\dot{u}_{k}\right)=w_{1,1} M_{\dot{\sigma}}^{n}\left(\begin{array}{l}
1 \\
0 \\
0 \\
0
\end{array}\right)
$$

We shall rather concentrate on the occurrences of $b^{+}$in $\dot{u}$. These will certainly take place, and thus $\dot{\sigma}^{n}\left(b^{+}\right)$will also occur for each $n$. But from the beginning to the end of such an occurrence the walk will travel

$$
w_{1,1} M_{\dot{\sigma}}^{n}\left(\begin{array}{c}
0  \tag{9}\\
0 \\
1 \\
0
\end{array}\right)=\beta^{n}
$$

steps (by (7)). Hence if $|\beta|>1$, then $\left(S_{N}(u)\right)$ is unbounded.
There remain 3 possibilities: $\beta=0$ or $\beta= \pm 1$.
In case $\beta=0, \dot{f}\left(\dot{\sigma}\left(b^{+}\right)\right)=\dot{f}\left(\dot{\sigma}\left(b^{-}\right)\right)=0$. But also $\dot{f}\left(\dot{\sigma}\left(a^{+} a^{-}\right)\right)=0$. Since symbols $a^{+}$and $a^{-}$alternate in $\dot{u}$, the sequence $\dot{u}$ has a decomposition in words from the set $V=\left\{\left(b^{+}\right)^{k}, a^{+}\left(b^{-}\right)^{k} a^{-}: k \geq 0\right\}$. Moreover, we may assume this set to be finite, since (by almost periodicity) the distance between the occurrence of two
$a$ 's in $u$ is bounded. Each word $v$ in the set $V$ has $\dot{f}(\dot{\sigma}(v))=0$. This clearly implies that $\left(\dot{S}_{N, \dot{f}}(\dot{u})\right.$ ), and hence $\left(S_{N}(u)\right)$ is bounded.
Now the case $\beta= \pm 1$. We see that $\beta$ is an eigenvalue of $M_{\dot{\sigma}}$ (with right eigenvector $(0,0,1,-1)^{T}$ ). Passing from $\sigma$ to $\sigma^{2}$ we may therefore assume (again by idempotency of $\bar{M}_{\sigma}$ ) that $\beta=+1$. In that case

$$
S_{\ell_{n}, \dot{f}}(\dot{u})=w_{1,1} M_{\dot{\sigma}}^{n}\left(\begin{array}{l}
1 \\
0 \\
0 \\
0
\end{array}\right)=1+(n-1) \alpha
$$

So if $\alpha \neq 0$, then the walk is unbounded. We are left with the case $\alpha=0, \beta=1$. Then we see from (6) that $w_{1,1}$ is a left eigenvector of $M_{\dot{\sigma}}$. Because of (8) this is a necessary and sufficient condition for $\dot{\sigma}$ to admit a representation by $\dot{f}$ in $\mathbf{R}$. (Cf. the remarks in the last paragraphs of the introduction). This implies that the walk is renormalizable in the following sense: if the 4 steps corresponding to the symbols $a^{+}, a^{-}, b^{+}$and $b^{-}$are replaced by the sequences of steps corresponding to $\dot{\sigma}\left(a^{+}\right), \dot{\sigma}\left(a^{-}\right), \dot{\sigma}\left(b^{+}\right)$, respectively $\dot{\sigma}\left(b^{-}\right)$, then this new walk is equal to the original walk. From this one deduces that if the $S_{N, \dot{f}}\left(\dot{\sigma}\left(a^{+}\right)\right), 1 \leq N \leq\left|\dot{\sigma}\left(a^{+}\right)\right|$crosses level 1 , then $S_{N, \dot{f}}\left(\dot{\sigma}^{2}\left(a^{+}\right)\right), 1 \leq N \leq\left|\dot{\sigma}^{2}\left(a^{+}\right)\right|$will cross level 2 . More generally $\left(S_{N, \dot{f}}\left(\dot{\sigma}^{2^{n}}\left(a^{+}\right)\right)\right.$) will cross level $n$ and the walk will be unbounded. So suppose $\left(S_{N, \dot{f}}(\dot{u})\right)_{N=0}^{\infty}$ remains between the levels -1 and +1 . Then $b^{2}$ cannot occur in $u$, as this would lead to three $+^{\prime} s$ or three $-' s$ in $\dot{u}$. Also since $\dot{f}\left(\dot{\sigma}\left(a^{+}\right)\right)=1$, and $\left(S_{N, \dot{f}}\left(\dot{\sigma}\left(a^{-}\right)\right)\right)_{N=0}^{\left|\dot{\sigma}\left(a^{-}\right)\right|-1}$ equals $S_{N, \dot{f}}\left(\dot{\sigma}\left(a^{+}\right)\right)_{N=0}^{\left|\dot{\sigma}\left(a^{+}\right)\right|-1}$ mirrored around zero, $a^{2}$ cannot occur, unless $\left(S_{N, \dot{f}}\left(\dot{\sigma}\left(a^{+}\right)\right)\right.$) stays between 0 and 1 , which would only be possible if $a^{+}$and $a^{-}$alternate in $\dot{\sigma}\left(a^{+}\right)$, what contradicts the primitivity of $M_{\dot{\sigma}}$. We have shown that $\sigma(a)$ contains neither $a^{2}$ nor $b^{2}$, but then $\sigma(a)=(a b)^{m} a$, where $m$ is even because $\sigma$ is of type III. Since the same arguments apply to $\sigma^{n}(a)$, and $\sigma(b)$ has to appear in some $\sigma^{n}(a), \sigma(b)$ will also neither contain $a^{2}$ nor $b^{2}$. Since $\sigma(a b)$ and $\sigma(b a)$ will occur, it follows likewise that the first and the last letter of $\sigma(b)$ are equal to $b$. Hence $\sigma(b)$ has the claimed form.
Example. Let $\tau$ be the Fibonacci substitution defined by $\tau(a)=a b, \tau(b)=a$. Then $\tau^{3}(a)=a b a a b, \tau^{3}(b)=a b a$, and $\sigma=\tau^{3}$ if of Type III. We have $\dot{\sigma}\left(a^{+}\right)=$ $a^{+} b^{-} a^{-} a^{+} b^{-}, \dot{\sigma}\left(b^{+}\right)=a^{+} b^{-} a^{-}$. hence $\alpha=-2$ and $\beta=-1$, so $\left(S_{N}(u)\right)$ is unbounded. (The substitution $\sigma$ given by $\sigma(a)=a b b, \sigma(b)=a b a b$ gives a (nonperiodic) example of a bounded walk.)

## 7. Type IV Substitutions

Here $\bar{M}_{\sigma}=\left(\begin{array}{ll}0 & 1 \\ 0 & 1\end{array}\right)$. Let us write $\nu=\sigma(a), \mu=\sigma(b)$. Then, because $\sigma$ is of Type IV, the Squaring Lemma implies

$$
\mu^{2} \sim \epsilon \text { and }(\mu \nu)^{2} \sim \epsilon
$$

So if we consider $\nu$ and $\mu$ as symbols, we have that the two groups

$$
G=<\bar{a}, \bar{b} \mid \bar{a}^{2}=(\bar{a} \bar{b})^{2}=\bar{\epsilon}>\text { and } H=<\bar{\mu}, \overline{\mu \nu} \mid \bar{\mu}^{2}=(\overline{\mu \nu})^{2}=\bar{\epsilon}>
$$

are isomorphic. Here $\bar{w}$ denotes the equivalence class of a word $w$ under the equivalence relation introduced in Section 4.
But the matrix $\left(\begin{array}{ll}0 & 1 \\ 0 & 1\end{array}\right)$ of $\sigma$ over $\{a, b\}$ transforms to the matrix $\left(\begin{array}{cc}1 & 0 \\ 1 & 0\end{array}\right)$ of $\hat{\sigma}$ over $\{\mu, \nu\}$, where the substitution $\hat{\sigma}$ on $\{\mu, \nu\}$ is defined as in Section 3. We then use the analysis of Type III, where at an occurrence of $\mu$ respectively $\nu$ we move $K:=$ $S_{\ell_{a}}(\sigma(a))$, respectively $L:=S_{\ell_{b}}(\sigma(b))$ steps $\left(\ell_{a}=|\sigma(a)|, \ell_{b}=|\sigma(b)|\right)$. Because $\bar{M}_{\hat{\sigma}}=\left(\begin{array}{l}1 \\ 1 \\ 1 \\ 0\end{array}\right)$, there are an odd number of $\nu$ 's in $\hat{\sigma}(\mu)$. But then the parameter $\alpha=s-t$ of the associated $\dot{\hat{\sigma}}$ matrix has to be odd, i.e., the renormalizable case $\alpha=0, \beta=1$ of the Type III analysis can not occur for these matrices. Furthermore, using the vector $w_{K, L}$ instead of $w_{1,1}$ in (9), we see from (7) that the walk is bounded iff $\beta=0$, which occurs iff $\dot{\hat{\sigma}}\left(\nu^{+}\right)$has an equal number of $\nu^{+}$ and $\nu^{-}$. Since $\dot{\hat{\sigma}}\left(\nu^{+}\right)$already has an equal number of $\mu^{+}$and $\mu^{-}$we finally obtain Proposition 3. If $\sigma$ if of Type IV, then $\left(S_{N}(u)\right)$ is bounded iff $\tau(b) \sim \epsilon$, where $\tau$ is the substitution obtained from $\sigma$ by interchanging a and $b$.
Example. Let $\sigma$ be defined by $\sigma(a)=a a b b, \sigma(b)=a b$. Then $\tau$ is given by $\tau(a)=b a, \tau(b)=b b a a$. Since $\tau(b) \nsim \epsilon, \sigma$ generates an unbounded oriented walk.
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[^0]:    Mots-clés : Substitutions, self-similarity, walks.

